

**Mini-lecture on Gödel's Constructible Universe**

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Forcing and Independence Proofs

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# Prerequisites

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## Reflection Theorems

Let  $ZFC^*$  be any sufficiently large **finite** fragment of the ZFC-axioms. Then:

- ▶ **Reflection #1.** For any  $a$  there is a set model  $M$  such that  $a \subseteq M$ ,  $M \models ZFC^*$  and  $|M| = \max(\aleph_0, |a|)$ .

(However:  $M$  is generally not transitive.)

- ▶ **Reflection #2.** For any  $a$  there is a set model  $M$  such that  $a \subseteq M$ ,  $M \models ZFC^*$  and  $M$  is transitive.

(However:  $|M|$  may be very large.)

↓  
 $M = V_\alpha$  is possible

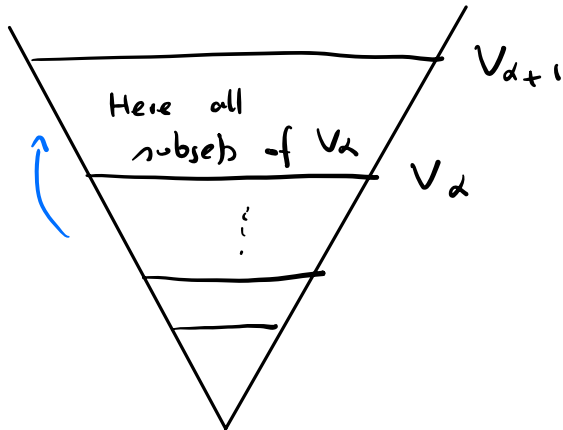
- ▶ **Reflection #3.** For any **transitive** set  $a$  there is a set model  $M$  such that  $a \subseteq M$ ,  $M \models ZFC^*$ ,  $M$  is transitive and  $|M| = \max(\aleph_0, |a|)$ .

E.g.:  $a = \{\aleph_{20}\}$

# Gödel's Constructible Universe $L$

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Recall  $V = \bigcup_{\alpha \in Ord} V_\alpha$



$$V_0 = \emptyset$$
$$V_{\alpha+1} = P(V_\alpha)$$
$$V_\lambda = \bigcup_{\alpha < \lambda} V_\alpha$$

Idea of  $L$ : Only use definable subsets at stages  $\alpha$ .

Definable  $\Rightarrow$  in FOL, lang. of set theory

## Definition

- ▶ A subset  $Y \subseteq X$  is **definable over  $X$**  if there is a formula  $\varphi$  and parameters  $a_1 \dots a_n \in X$  such that

$$Y = \{y \in X \mid X \models \varphi(y, a_1 \dots a_n)\}.$$

- ▶  $\text{Def}(X) = \{Y \subseteq X \mid Y \text{ is definable over } X\}.$

→ Important that  $X$  is a set, and so you can formally express " $\exists \varphi$  st.  $X \models \varphi \dots$ "

## Definition

- ▶  $L_0 := \emptyset.$
- ▶  $L_{\alpha+1} := \text{Def}(L_\alpha).$
- ▶  $L_\lambda := \bigcup_{\alpha < \lambda} L_\alpha$  (for  $\lambda$  limit).

$$L := \bigcup_{\alpha \in \text{Ord}} L_\alpha$$

Let's see how  $V$  differs from  $L$ .

- $L_0 = V_0 = \emptyset$
- $L_n = V_n$  for  $n < \omega$  (because fin. sets are definable)
- $L_\omega = V_\omega$
- $L_{\omega+1} \stackrel{?}{=} V_{\omega+1}$

NB:  $|V_{\omega+1}| = 2^{\aleph_0}$

$|L_{\omega+1}| = \aleph_0$ , because there are only cfbly many formulas

$$|V_{\omega+2}| = 2^{(2^{\aleph_0})} \quad \text{while} \quad |L_{\omega+2}| = \aleph_0$$

Lemma: For  $\alpha > \omega$ ,  $|L_\alpha| = |\alpha|$ .

Proof: Induction +  $|\text{Def}(x)| = |x|$  if  $x$  is finite.

Facts: •  $L_\alpha$  is transitive

•  $\alpha < \beta \Rightarrow L_\alpha \in L_\beta$

•  $\text{Ord} \cap L_\alpha = \alpha$  (i.e.  $\beta \in L_\alpha \Leftrightarrow \beta < \alpha$ )

Define:  $x \in \mathbb{L} \Rightarrow \rho_L(x) := \text{least } \alpha \text{ s.t. } x \in L_{\alpha+1}$ .

## Theorem

$L \models ZF$

⋮

• Power Set: let  $x \in L_\beta$ , so  $x \in L_\alpha$ .

Need to find  $y \in L$  s.t.

$$L \models y = P(x)$$

$$L \models \forall z (z \in y \leftrightarrow z \in x)$$

For every  $z \in x$ , if  $z \in L$  then let

$$\alpha_z := \beta_L(z) \quad (\alpha_z > \alpha)$$

Let  $\beta := \sup \{ \alpha_z \mid z \in x \text{ and } z \in L \}$

$$P(x)^L := \{ z \mid z \in x \wedge z \in L \} = P(x) \cap L \subseteq L_{\beta+1}$$

Then: use Comprehension in  $L$ .

• For Comprehension: you need Reflection to shift  
from " $L \models \varphi$ " to " $L_\delta \models \varphi$ ".

⋮

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## $L \models AC$ (without assuming AC)

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In fact  $L \models$  Global Choice, i.e. there is a (class) wellorder  $<_L$  of all of  $L$ .

Proof: Inductively. Assume  $L_\alpha$  is wellordered by  $<_\alpha$ . It follows  $L_\alpha^{<\omega}$  are wellordered by  $<_\alpha^{lex}$ .

Now: define  $<_{\alpha+1}$  on  $L_{\alpha+1}$ :

• Given  $x, y \in L_{\alpha+1}$ :

(1) If  $x, y \in L_\alpha \Rightarrow (x <_{\alpha+1} y \Leftrightarrow x <_\alpha y)$

(2) If  $x \in L_\alpha, y \in L_{\alpha+1} \setminus L_\alpha \Rightarrow x <_{\alpha+1} y$

(3) If  $x, y \in L_{\alpha+1} \setminus L_\alpha \Rightarrow$

Definable by some  $\varphi, \psi$ , and  $a_1, \dots, a_n \in L_\alpha$ .

$x <_{\alpha+1} y \Leftrightarrow$  the smallest  $\ulcorner \varphi \urcorner$  and  $<_\alpha^{lex}$ -least  $\langle a_1, \dots, a_n \rangle$  s.t.  $x$  is def. by  $\varphi(\dots, a_1, \dots, a_n)$

is  $<_\alpha^{lex}$ -smaller than

the smallest  $\ulcorner \psi \urcorner$  and  $<_\alpha^{lex}$ -least  $\langle b_1, \dots, b_k \rangle$  s.t.

$y$  is def. by  $\psi(\dots, b_1, \dots, b_k)$ .

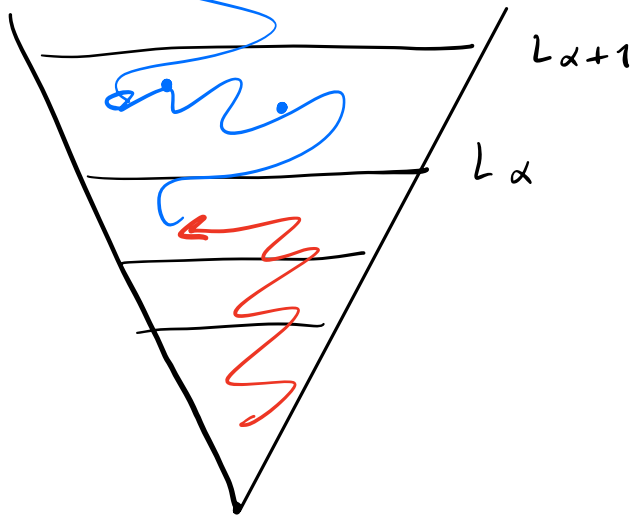
$$x = \{z \mid L_\alpha \models \varphi(z, a_1, \dots, a_n)\}$$

$$y = \{z \mid L_\alpha \models \psi(z, b_1, \dots, b_k)\}$$

$$\langle \ulcorner \varphi \urcorner, a_1, \dots, a_n \rangle$$

$$\langle \ulcorner \psi \urcorner, b_1, \dots, b_k \rangle$$

which comes first?





## Absoluteness and the Axiom of Constructibility

### Theorem

$L$  is absolute for transitive models.

More specifically, the function  $\alpha \mapsto L_\alpha$  is absolute for transitive models.

Corollary: For any transitive  $M$ ,  $(M \models ZF^*)$   
 $\alpha \in M \iff L_\alpha \in M$

Proof: Sketchy: First-order logic  $(M \models \varphi)$   
is defined by recursion  
using basic operations  
 $\Rightarrow$  absolute.  
 $L_\alpha :=$  defined by recursion using  
 $\text{Def}(x) \Rightarrow$  absolute.

Check:  $\text{Def}(x)^M = \text{Def}(x)$   
 $\vdots$   
 $\boxtimes$

### Definition

The statement " $\forall x \exists \alpha \in \text{Ord} (x \in L_\alpha)$ " is called the **Axiom of Constructibility**, and usually abbreviated by

$$V = L.$$

### Corollary 1

$$L \models (V = L)$$

Proof:  $\forall x \in L \exists \alpha \in \text{Ord} \quad x \in L_\alpha$  (by def).

$$\Leftrightarrow \forall x \in L \exists \alpha \in \text{Ord} \quad (x \in L_\alpha)^L$$
$$\Leftrightarrow \forall x \in L \exists \alpha \in \text{Ord} \cap L \quad (x \in L_\alpha)^L$$
$$\Leftrightarrow [V = L]^L$$

### Corollary 2

If  $M$  is any transitive class model of  $ZF^*$  with  $\text{Ord} \subseteq M$ , then  $L \subseteq M$ .

Proof:  $\alpha \in M \Rightarrow L_\alpha \in M \Rightarrow L_\alpha \in M \Rightarrow L \in M \quad \square$

This means:  $L$  is the minimal model of  $ZFC$ .

### Corollary 3

If  $M$  is any transitive class model of  $ZF^* + V = L$  with  $\text{Ord} \subseteq M$ , then  $L = M$ .

Proof:  $M \models V = L$ , then  $M \models \forall x \exists \alpha (x \in L_\alpha)$

$$\text{So } \forall x \in M \exists \alpha \in \text{Ord} \cap M \quad (x \in L_\alpha)^M$$

$\xrightarrow{\text{abs}}$

$$\forall x \in M \exists \alpha \in \text{Ord} \quad (x \in L_\alpha)$$
$$\forall x \in M : x \in L$$
$$\text{So } M \subseteq L$$

NB:  $ZF^*$  has to be enough to prove absoluteness of  $L_\alpha$ 's.

## Miniature versions

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### Definition

If  $M$  is a transitive set model, then the **height** of  $M$ , denoted by  $o(M)$ , is defined as the least ordinal not in  $M$  (alternatively,  $o(M) := \text{Ord} \cap M$ ).

### Corollary 4

$L_\delta \models (\mathbf{V} = \mathbf{L})$  for all limit  $\delta$ .

Just as before.

### Corollary 5

If  $M$  is any set model of  $\text{ZF}^*$  with height  $\delta$ , then  $L_\delta \subseteq M$ .

$$\alpha \in M \Rightarrow \alpha < \delta \Rightarrow L_\alpha \in M. \quad \text{So } L_\alpha \in M, \text{ so}$$

$$L_\delta = \bigcup_{\alpha < \delta} L_\alpha \in M$$

### Corollary 6

If  $M$  is any set model of  $\text{ZF} + \mathbf{V} = \mathbf{L}$  with height  $\delta$ , then  $M = L_\delta$ .

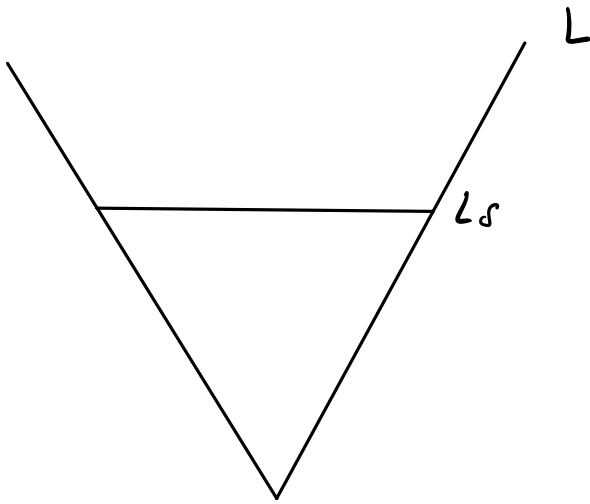
$$M \models (\mathbf{V} = \mathbf{L})$$

$$M \models \forall x \exists \alpha \ x \in L_\alpha$$

$$\forall x \in M \exists \alpha \in \text{Ord} \cap M \ (x \in L_\alpha)^M$$

$$\forall x \in M \exists \alpha < \delta \ (x \in L_\alpha)$$

$$\text{So } x \in M \Rightarrow x \in \bigcup_{\alpha < \delta} L_\alpha = L_\delta \quad \text{so } M \subseteq L_\delta.$$



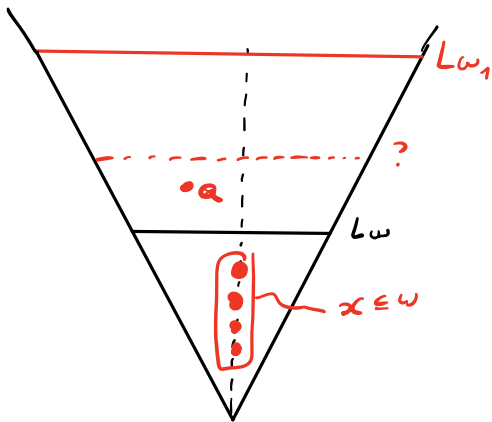
Only  $M \neq 2F + (v=L)$   
are  $L\delta$ 's  
where  $\delta = o(M)$

# Condensation and GCH

## Condensation Lemma

Suppose  $V = L$ . Any  $x \subseteq \omega$  is contained in  $L_{\omega_1}$ .

More generally, any  $x \subseteq \kappa$  is contained in  $L_{\kappa^+}$ .



Proof: Let  $x \subseteq \omega$  and let  $a = \omega \cup \{x\}$  (to make it transitive)

By Reflection #3: there is  $M \models ZF^* + V=L$ , s.t.

- $a \in M$
- $|M| = \max(|a|, \aleph_0) = \aleph_0$
- $M$  transitive

But then:  $M = L_\delta$ , for  $\delta = o(M)$ .

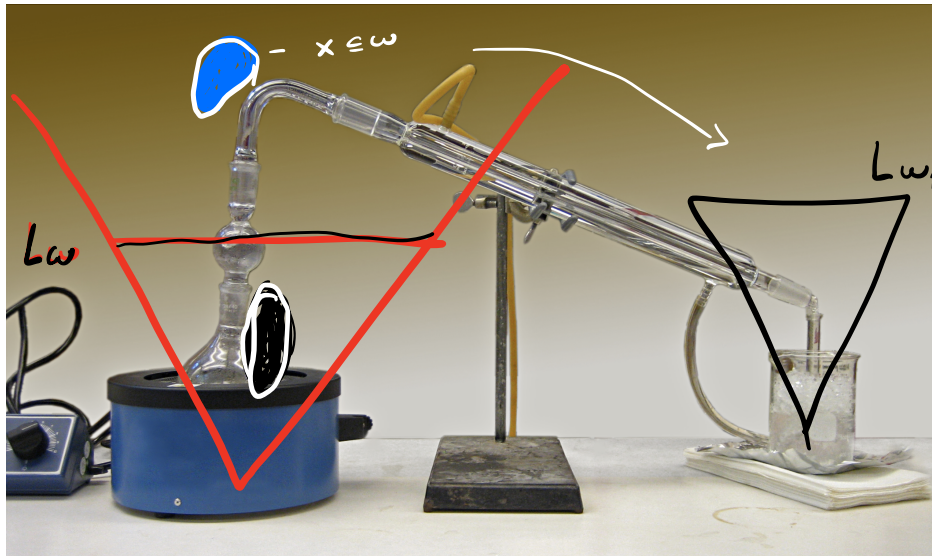
What's  $o(M) < \omega_1$  (a countable ordinal)

( $o(M) = M \cap \text{Ord}$  and  $|M| = \omega$ )

So  $\delta < \omega_1$ .

So  $a \in L_\delta \subseteq L_{\omega_1}$ . So  $x \in a \in L_{\omega_1}$ .

So  $x \in L_{\omega_1}$ . □



Corollary

$L \models \text{GCH}$

Proof: Assume  $V=L$ .

Since any  $x \in \kappa$  is in  $L_{\kappa^+}$ , so

$$P(\kappa) \subseteq L_{\kappa^+}$$

$$2^\kappa = |P(\kappa)| \leq |L_{\kappa^+}| = |\kappa^+| = \kappa^+ \leq 2^\kappa$$

□